

# Tensor Numerical Methods

Fundamental principles and application to TEI calculations

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# 0. Motivation

Consider a real rectangular matrix  $\mathbf{A} \in \mathbb{R}^{n \times n}$  and a vector  $\mathbf{v} \in \mathbb{R}^n$ .  
Computation of matrix vector product

$$\mathbf{A}\mathbf{v} = \mathbf{x} \in \mathbb{R}^m \quad (1)$$

costs  $n^2$  operations.

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If we know the **S**ingular **V**alue **D**ecomposition of  $\mathbf{A}$ :

$$\mathbf{A} = \mathbf{U}\mathbf{S}\mathbf{V}^T \quad (2)$$

where  $\mathbf{U} \in \mathbb{R}^{m \times r}$ ,  $\mathbf{V} \in \mathbb{R}^{n \times r}$  are orthonormal matrices,  $\mathbf{S} \in \mathbb{R}^{r \times r}$  is diagonal matrix and  $r = \text{rank}(\mathbf{A})$ , we can rewrite (1) as

$$\mathbf{U}\mathbf{S}\mathbf{V}^T\mathbf{v} = \mathbf{x}. \quad (3)$$

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For simplicity consider an orthogonal (generally not orthonormal) matrix

$\tilde{\mathbf{V}} = (\mathbf{S}\mathbf{V}^T)^T$ , so

$$\mathbf{A} = \mathbf{U}\tilde{\mathbf{V}}^T. \quad (4)$$

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However, rank  $r$  does not have to be low.

Solution = Truncated SVD Let  $\mathbf{USV}^T$  be the SVD decomposition of  $\mathbf{A}$ . Without loss of generality consider a non-increasing sequence of singular values

$$\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_n \geq 0, \quad (5)$$

so  $\mathbf{S} = \text{diag}([\sigma_1, \sigma_2, \dots, \sigma_n])$ . Set

$$\mathbf{S}_k = \text{diag}([\sigma_1, \sigma_2, \dots, \sigma_k, 0, \dots, 0]) \in \mathbb{R}^{n \times n} \quad (6)$$

$$\bar{\mathbf{S}}_k = \text{diag}([\sigma_1, \sigma_2, \dots, \sigma_k]) \in \mathbb{R}^{k \times k} \quad (7)$$

and matrices  $\bar{\mathbf{U}}_k, \bar{\mathbf{V}}_k$  as first  $k$  columns of matrices  $\mathbf{U}, \mathbf{V}$ . Define

$$\mathbf{A}_k := \mathbf{US}_k\mathbf{V}^T = \bar{\mathbf{U}}_k\bar{\mathbf{S}}_k\bar{\mathbf{V}}_k^T. \quad (8)$$

$$\mathbf{A}_k := \mathbf{U}\mathbf{S}_k\mathbf{V}^T = \bar{\mathbf{U}}_k\bar{\mathbf{S}}_k\bar{\mathbf{V}}_k^T.$$

Then  $\mathbf{A}_k$  is the best rank- $k$  approximation of  $\mathbf{A}$ , i.e.

$$\|\mathbf{A} - \mathbf{B}\| \geq \|\mathbf{A} - \mathbf{A}_k\|, \text{ rank}(\mathbf{B}) \leq k. \quad (9)$$

(Eckart-Young Theorem).

Error of approximation:

$$\|\mathbf{A} - \mathbf{A}_k\|_F = \sqrt{\sum_{i=k+1}^n \sigma_i^2} \quad (10)$$

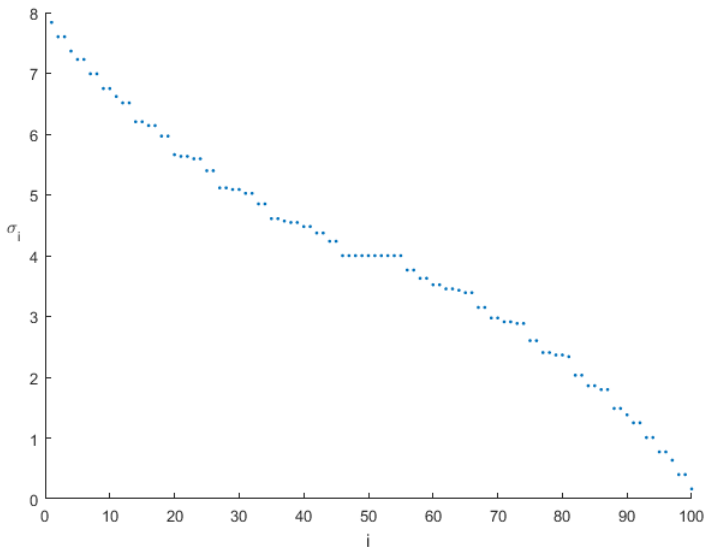
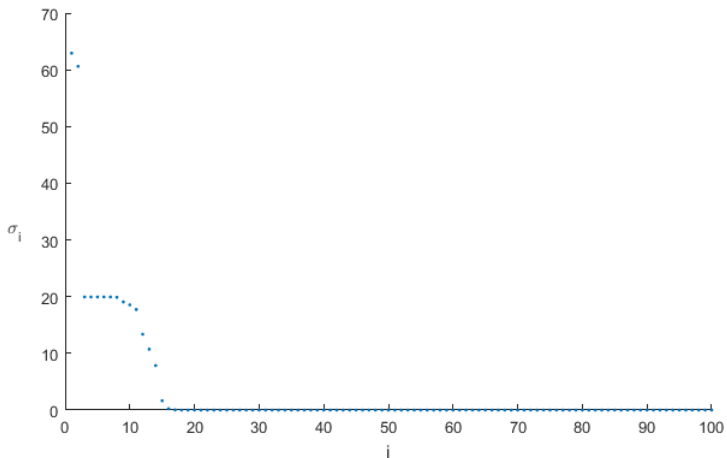
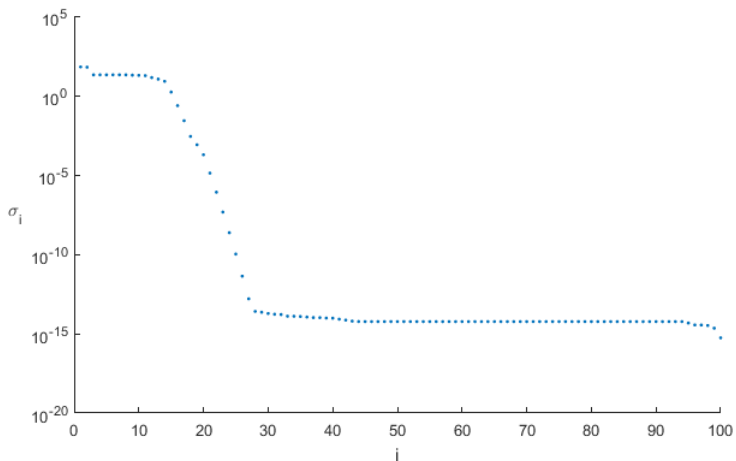


Figure: Singular values of discretized 2D laplacian (matrix  $100 \times 100$ ).



**Figure:** Singular values of discretized function  $f(x, y) = \sin(x + y) \cos(xy) - y^2 \cos(x)$  on interval  $[0, 2\pi] \times [0, 2\pi]$  (matrix  $100 \times 100$ ).



**Figure:** Singular values of discretized function  $f(x, y) = \sin(x + y) \cos(xy) - y^2 \cos(x)$  on interval  $[0, 2\pi] \times [0, 2\pi]$  (matrix  $100 \times 100$ ). Logarithmic scale.

## Definition

A tensor of order  $d$  ( $N$ - $d$  tensor) is a multidimensional array over a  $d$ -tuple index set,

$$\mathbf{A} \in \mathbb{R}^{n_1 \times n_2 \times \dots \times n_d}, [\mathbf{A}]_{i_1 i_2 \dots i_d} = a_{i_1 i_2 \dots i_d} \in \mathbb{R}. \quad (11)$$



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An  $N$ - $d$  tensor = an element of a tensor-product Hilbert space (finite dimensional)

$$\mathbb{W}_{\mathbf{n}} \equiv \mathbb{W}_{\mathbf{n},d} = \bigotimes_{\ell=1}^d X_{\ell} \quad (12)$$

where

$$X_{\ell} = \mathbb{R}^{n_{\ell}}, \mathbf{n} = (n_1, \dots, n_d). \quad (13)$$

# Multilinear algebra operations

## Scalar product

Scalar (inner) product is defined as

$$\langle \mathbf{A}, \mathbf{B} \rangle := \sum_{i \in \mathcal{I}} [\mathbf{A}]_i [\mathbf{B}]_i, \quad \mathbf{A}, \mathbf{B} \in \mathbb{W}_n. \quad (14)$$

Induced norm:

$$\|\mathbf{A}\| := \sqrt{\langle \mathbf{A}, \mathbf{A} \rangle} = \sqrt{\sum_{i \in \mathcal{I}} [\mathbf{A}]_i^2}. \quad (15)$$

(Analogy to Frobenius norm)

# Multilinear algebra operations

Hadamard product, Discrete convolution

Hadamard (entrywise) product  $\odot : \mathbb{W}_n \times \mathbb{W}_n \rightarrow \mathbb{W}_n$  is defined as

$$[\mathbf{A} \odot \mathbf{B}]_i = [\mathbf{A}]_i [\mathbf{B}]_i, \mathbf{i} \in \mathcal{I}. \quad (16)$$

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For given tensors  $\mathbf{A}$  and  $\mathbf{B}$  the discrete convolution can be defined entrywise as

$$[\mathbf{A} * \mathbf{B}]_j = \left[ \sum_{\mathbf{k} \in \mathcal{I}} [\mathbf{A}]_{\mathbf{k}} [\mathbf{B}]_{j-\mathbf{k}} \right]_j, \quad (17)$$

$\underbrace{\{1, 2, \dots, 2n_1 - 1\} \times \dots \times \{1, 2, \dots, 2n_d - 1\}}_{=:\mathcal{J}}$

# Multilinear algebra operations

## Computational cost

Operations are very expensive for large  $d$ :

- scalar product:  $\mathcal{O}(n^d)$
- Hadamard product:  $\mathcal{O}(n^d)$
- discrete convolution:  $\mathcal{O}(n^{2d})$  ( $n^d \log n$  if we use the fast FFT algorithm)

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**"The Curse of Dimensionality".**

# Contracted product

Given a tensor  $\mathbf{A} \in \mathbb{R}^{n_1, \dots, n_{\ell-1}, n_\ell, n_{\ell+1}, \dots, n_d}$  and a matrix  $\mathbf{B} \in \mathbb{R}^{n_\ell, m}$ , we define a contracted product along mode  $\ell$  as

$$\mathbb{R}^{n_1, \dots, n_{\ell-1}, m, n_{\ell+1}, \dots, n_d} \ni \mathbf{C} = \mathbf{A} \times_\ell \mathbf{B} \quad (18)$$

where each entry of the tensor  $\mathbf{C}$  may be written as

$$c_{i_1, \dots, i_{\ell-1}, j_\ell, i_{\ell+1}, \dots, i_d} = \sum_{k=1}^{n_\ell} a_{i_1, \dots, i_{\ell-1}, k, i_{\ell+1}, \dots, i_d} b_{k, j_\ell}, \quad j_\ell = 1, \dots, m. \quad (19)$$



# Tensor reshaping

## Vectorization of tensor

Consider a matrix  $\mathbf{A} \in \mathbb{R}^{m \times n}$ . The vector representation  $\text{vec}(\mathbf{A}) \in \mathbb{R}^{mn}$  = putting the matrix columns (ordered) into the vector

$$\text{vec}(\mathbf{A}) = [a_{11}, \dots, a_{m1}, a_{12}, \dots, a_{m2}, \dots, a_{1n}, \dots, a_{mn}]^T, \quad (20)$$

or

$$\text{vec}(\mathbf{A}) = \begin{bmatrix} [a(:, 1)] \\ [a(:, 2)] \\ \vdots \\ [a(:, n)] \end{bmatrix}. \quad (21)$$

# Tensor reshaping

## Vectorization of tensor

In case of arbitrary tensor  $\mathbf{B} \in \mathbb{R}^{n_1 \times \dots \times n_d}$  we can define it recursively:

$$\text{vec}(\mathbf{B}) = \begin{bmatrix} \text{vec}([b(:, \dots, :, 1)]) \\ \text{vec}([b(:, \dots, :, 2)]) \\ \vdots \\ \text{vec}([b(:, \dots, :, n)]) \end{bmatrix}. \quad (22)$$

"Stacking" of 1-mode fibers of  $\mathbf{B}$ .

Index mapping:

$$j = 1 + \sum_{k=1}^d (i_k - 1) \prod_{\ell=1}^{k-1} n_{\ell}. \quad (23)$$

# Tensor reshaping

## Matrix unfolding

Matrix unfolding of a given tensor  $\mathbf{B} \in \mathbb{R}^{n_1 \times \dots \times n_d}$  along mode  $\ell$  is a matrix  $\mathbf{B}_{(\ell)} \in \mathbb{R}^{n_\ell \times (n_1 \dots n_{\ell-1} n_{\ell+1} \dots n_d)}$  which columns are ordered  $\ell$ -mode fibers of  $\mathbf{B}$ .

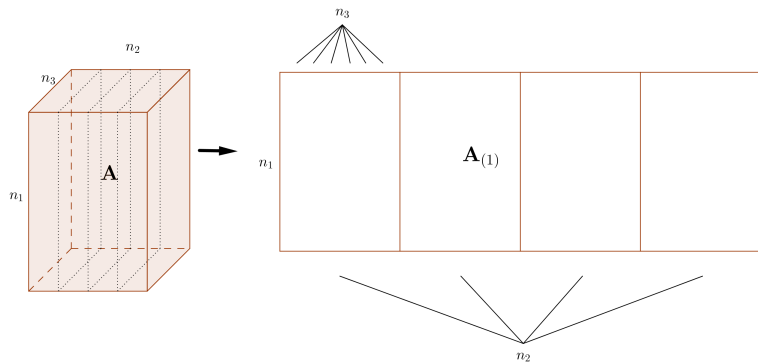


Figure: Matrix unfolding of 3-d tensor  $\mathbf{A}$  along mode 1

## 2. Problem setting

# Calculation of TEI

To solve discretized Hartree-Fock equations (e.g. using the Direct Inversion of the Iterative Subspace (DIIS) or Inexact Restoration method (IRM)) we have to precompute so-called two-electron integrals:

$$b_{ijkl} = \int_{\mathbb{R}^3} \int_{\mathbb{R}^3} \frac{g_i(\mathbf{x}) g_j(\mathbf{x}) g_k(\mathbf{y}) g_l(\mathbf{y})}{\|\mathbf{x} - \mathbf{y}\|} d^3\mathbf{y} d^3\mathbf{x}, \quad (24)$$

where basis functions  $g_* : \mathbb{R}^3 \rightarrow \mathbb{R}$  satisfy  $|g_*(\mathbf{x})| \rightarrow 0$  when  $\|\mathbf{x}\| \rightarrow \infty$ .

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For gaussian basis functions

$$g(\mathbf{x}) = Np(\mathbf{x})e^{-\alpha\|\mathbf{x}\|^2}, \quad (25)$$

where  $p(\mathbf{x}) = p(x_1, x_2, x_3) = x_1^k x_2^l x_3^m$ ,  $k, l, m \in \mathbb{N}_0$ , an analytical solution is known.

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Consider a finite box (e.g. a cube) with an equidistant discretization with step size  $h$ . So each basis function is approximated by 3-d tensor

$$g_*(\mathbf{x}) \approx \mathbf{G}_* \in \mathbb{R}^{n^{\otimes 3}}. \quad (26)$$

If we rewrite the TEI as

$$b_{ijkl} = \int_{\mathbb{R}^3} g_i(\mathbf{x}) g_j(\mathbf{x}) \int_{\mathbb{R}^3} g_k(\mathbf{y}) g_l(\mathbf{y}) \frac{1}{\|\mathbf{x} - \mathbf{y}\|} d^3\mathbf{y} d^3\mathbf{x}, \quad (27)$$

we can understand the inner integral as function of  $\mathbf{x}$ .



The discretized problem:

$$b_{ijkl} \approx h^6 \langle \mathbf{G}_i \odot \mathbf{G}_j, \mathcal{P}_{\mathcal{N}} * (\mathbf{G}_k \odot \mathbf{G}_l) \rangle, \quad (28)$$

where  $\mathcal{P}_{\mathcal{N}} \in \mathbb{R}^{(2n-1)^{\otimes 3}}$  is discretized Newton convolving kernel  $\frac{1}{\|\mathbf{y}\|}$ . To resolve singularities we have to consider very fine grid - in order of  $10^5$  -  $10^6$  in one dimension.

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Solution - low rank tensor approximations.

### 3. Tensor formats

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- Orthonormality, no diagonal core - **Tucker format**



# Canonical format

A so-called rank-1 tensor  $\mathbf{V} \in \mathbb{W}_n$  can be written as a tensor product

$$\mathbf{V} = \mathbf{v}^{(1)} \otimes \dots \otimes \mathbf{v}^{(d)} = \bigotimes_{\ell=1}^d \mathbf{v}^{(\ell)}, \quad \mathbf{v}^{(\ell)} \in \mathbb{R}^{n_\ell}. \quad (29)$$

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Each tensor  $\mathbf{A} \in \mathbb{W}_n$  may be written as

$$\mathbf{A} = \sum_{k=1}^R \mathbf{a}_k^{(1)} \otimes \dots \otimes \mathbf{a}_k^{(d)}, \quad \mathbf{a}_k^{(\ell)} \in \mathbb{R}^{n_\ell}. \quad (30)$$

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That is so called Canonical  $R$ -term representation (we write  $\mathbf{A} \in \mathcal{C}_R(\mathbb{W}_n)$  or  $\mathbf{A} \in \mathcal{C}_{R,n}$  which means that  $\text{rank}(\mathbf{A}) \leq R$ ). Minimal  $R \in \mathbb{N}$  - canonical rank. It satisfies

$$1 \leq R \leq n^{d-1}. \quad (31)$$

$i$ -th entry of  $\mathbf{A}$ , we simply evaluate

$$[\mathbf{A}]_i = \sum_{k=1}^R [\mathbf{a}_k^{(1)}]_{i_1} \cdot \dots \cdot [\mathbf{a}_k^{(d)}]_{i_d}. \quad (32)$$

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Sometimes we require vectors  $\mathbf{a}_k^{(\ell)}$  to be normalized. Then we substitute

$$\mathbf{a}_k^{(\ell)} = c_k^{(\ell)} \mathbf{u}_k^{(\ell)}, \quad c_k^{(\ell)} \in \mathbb{R}, \quad \|\mathbf{u}_k^{(\ell)}\| = 1. \quad (33)$$

So the equivalent normalized canonical representation is written as

$$\mathbf{A} = \sum_{k=1}^R c_k \mathbf{u}_k^{(1)} \otimes \dots \otimes \mathbf{u}_k^{(d)}, \quad (34)$$

$c_k$  coefficients are similar to singular values - possible rank reduction.

# Canonical format $\mathcal{C}_{R,n}$

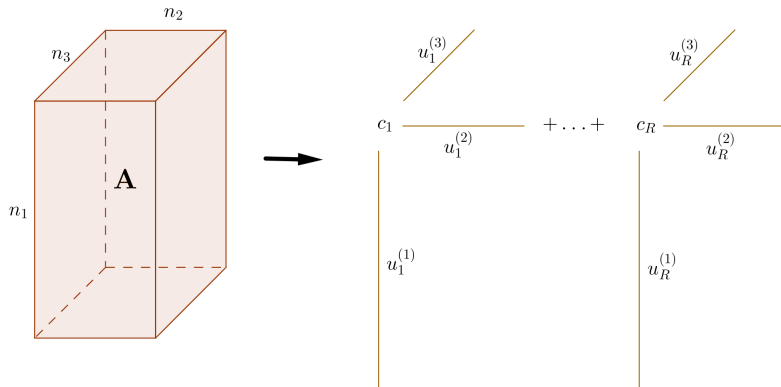


Figure: Canonical representation of tensor  $\mathbf{A}$

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pros and cons

- + Storage cost  $dRn$
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- + Storage cost  $dRn$
- + Cheap tensor operations
  - No algorithm to get canonical rank (set of "max  $R$ -rank tensors" not closed)
  - No algorithm to find the best rank- $R$  approximation (does not have to exist)



# Tucker format

Given the rank parameter  $\mathbf{r} = (r_1, \dots, r_d)$  we define the orthogonal Tucker format as

$$\mathbf{A} := \sum_{\mathbf{k}=\mathbf{1}}^{\mathbf{r}} b_{k_1, \dots, k_d} \mathbf{v}_{k_1}^{(1)} \otimes \dots \otimes \mathbf{v}_{k_d}^{(d)} \quad (35)$$

where  $\mathbf{k} = (k_1, \dots, k_d)$  and  $\forall \ell \in \{1, \dots, d\}$  the set of vectors  $\mathbf{v}_{k_\ell}^{(\ell)}$ ,  $1 \leq k_\ell \leq r_\ell$  is orthonormal. Coefficients  $b_{k_1, \dots, k_d}$  form a tensor  $\mathbf{B}$  of order  $d$ . We write  $\mathbf{A} \in \mathcal{T}_{\mathbf{r}}(\mathbb{W}_{\mathbf{n}})$  or  $\mathbf{A} \in \mathcal{T}_{\mathbf{r}, \mathbf{n}}$ .

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$i$ -th element of the tensor  $\mathbf{A}_{(\mathbf{r})}$ :

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# Tucker format $\mathcal{T}_{r,n}$

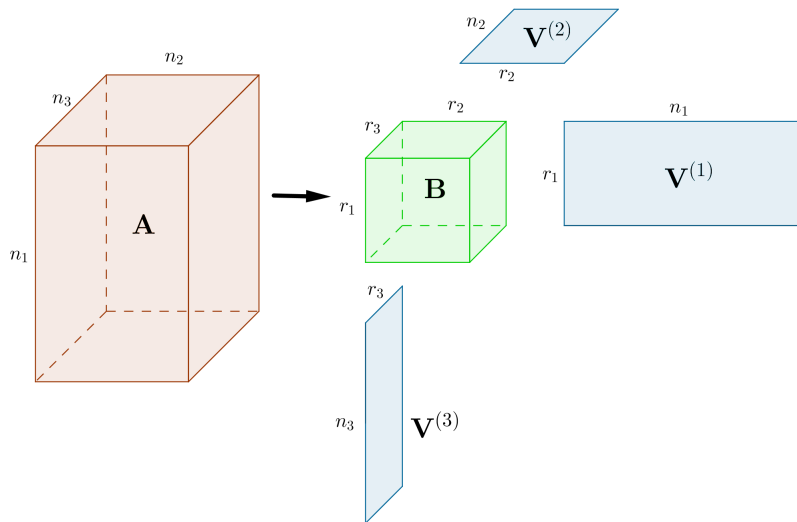


Figure: Tensor  $\mathbf{A}$  in Tucker format

# Tucker format $\mathcal{T}_{r,n}$

## Existence of a Tucker decomposition

### Theorem

(higher order SVD - HOSVD)

Each tensor  $\mathbf{A} \in \mathbb{R}^{n_1 \times \dots \times n_d}$  can be written as the product

$$\mathbf{A} = \mathbf{S} \times_1 \mathbf{V}^{(1)} \times_2 \dots \times_d \mathbf{V}^{(d)}, \quad (38)$$

where  $\mathbb{R}^{n_\ell \times n_\ell} \ni \mathbf{V}^{(\ell)} = [\mathbf{v}_1^{(\ell)} \mathbf{v}_2^{(\ell)} \dots \mathbf{v}_{n_\ell}^{(\ell)}]$  are orthonormal matrices ( $(\mathbf{V}^{(\ell)})^T \mathbf{V}^{(\ell)} = \mathbf{I} \in \mathbb{R}^{n_\ell \times n_\ell}$ ) and the tensor  $\mathbf{S}$  has the same size as  $\mathbf{A}$ . For fixed index  $\ell$  sub-tensors of  $\mathbf{S}$ ,  $\mathbf{S}_{i_\ell=\alpha}$ , have the following properties

- 1  $\langle \mathbf{S}_{i_\ell=\alpha}, \mathbf{S}_{i_\ell=\beta} \rangle = 0$  when  $\alpha \neq \beta$  (all-orthogonality),
- 2  $\|\mathbf{S}_{i_\ell=1}\| \geq \|\mathbf{S}_{i_\ell=2}\| \geq \dots \geq \|\mathbf{S}_{i_\ell=n_\ell}\| \geq 0$  (ordering).
- 3  $\mathbf{S} = \mathbf{A} \times_1 (\mathbf{V}^{(1)})^T \times_2 \dots \times_d (\mathbf{V}^{(d)})^T$ , ( $\mathbf{S}$  is the projection of  $\mathbf{A}$  onto the tensor basis formed by orthonormal matrices  $\mathbf{V}^{(\ell)}$ )

Frobenius norms  $\|\mathbf{S}_{i_\ell=j}\|$  are equal to the singular values  $\sigma_j^{(\ell)}$  of the matrix unfolding  $\mathbf{A}_{(\ell)}$  ( $\ell$ -mode singular values). Vectors  $\mathbf{v}_j^{(\ell)}$  are corresponding left singular vectors of  $\mathbf{A}_{(\ell)}$ .

### Algorithm

(HOSVD)

- 1 Given a tensor  $\mathbf{A}$ , for each  $\ell = 1, \dots, d$  construct the matrix unfolding  $\mathbf{A}_{(\ell)}$ .
- 2 For each unfolding find an SVD decomposition, i.e.  
 $\mathbf{A}_{(\ell)} = \mathbf{V}^{(\ell)} \Sigma^{(\ell)} (\mathbf{W}^{(\ell)})^T$ . Store the matrices  $\mathbf{V}^{(\ell)}$ .
- 3 Compute the core tensor as  $\mathbf{S} = \mathbf{A} \times_1 (\mathbf{V}^{(1)})^T \times_2 \dots \times_d (\mathbf{V}^{(d)})^T$ .

### Theorem

*(Truncated HOSVD)*

Let the HOSVD of the tensor  $\mathbf{A}$  is given and let the rank of matrix unfolding  $\mathbf{A}_{(\ell)}$  be equal to  $R_\ell$  for  $\ell = 1, \dots, d$ . Let's denote by  $\tilde{\mathbf{A}}$  the tensor which is obtained from  $\mathbf{A}$  by discarding the smallest  $\ell$ -mode singular values  $\sigma_{r_\ell+1}^{(\ell)}, \sigma_{r_\ell+2}^{(\ell)}, \dots, \sigma_{R_\ell}^{(\ell)}$  for given values of  $r_\ell$  ( $\ell = 1, \dots, d$ ). Then the error of the approximation satisfies

$$\|\mathbf{A} - \tilde{\mathbf{A}}\|^2 \leq \sum_{\ell=1}^d \sum_{i_\ell=r_\ell+1}^{R_\ell} (\sigma_{i_\ell}^{(\ell)})^2. \quad (39)$$

# Tucker format

truncated HOSVD

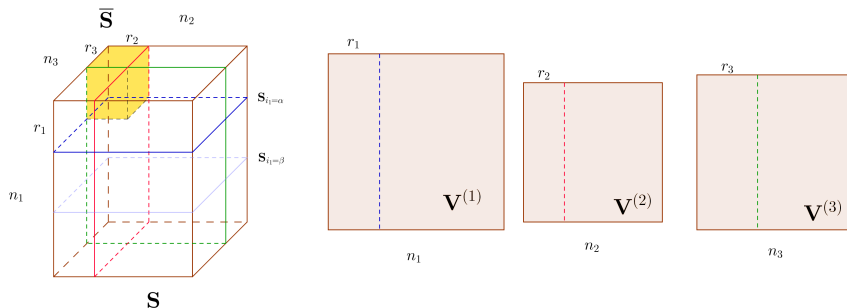


Figure: Mechanism of truncated HOSVD



# Tucker format

## pros and cons

- + Storage cost  $r^d + drn$
- + Cheap tensor operations (more expensive than in case of can. format)
- + Quite robust algorithms to find the best rank- $\mathbf{r}$  approximation
- + Efficient especially for 3D tensors

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- + Efficient especially for 3D tensors
  - Expensive for larger  $d$

## Definition

Given rank parameters  $R, \mathbf{r}$  we denote by  $\mathcal{T}_{C_{R,\mathbf{r}}}$  a subclass of tensors written in Tucker format  $\mathcal{T}_{\mathbf{r},\mathbf{n}}$  whose core tensor  $\mathbf{B}$  is represented in the canonical format.

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A tensor in two-level format can be written as

$$\mathbf{A} = \left( \sum_{k=1}^R \xi_k \mathbf{u}_k^{(1)} \otimes \dots \otimes \mathbf{u}_k^{(d)} \right) \times_1 \mathbf{V}^{(1)} \times_2 \dots \times_d \mathbf{V}^{(d)}. \quad (40)$$

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- storage cost  $dRr + drn$

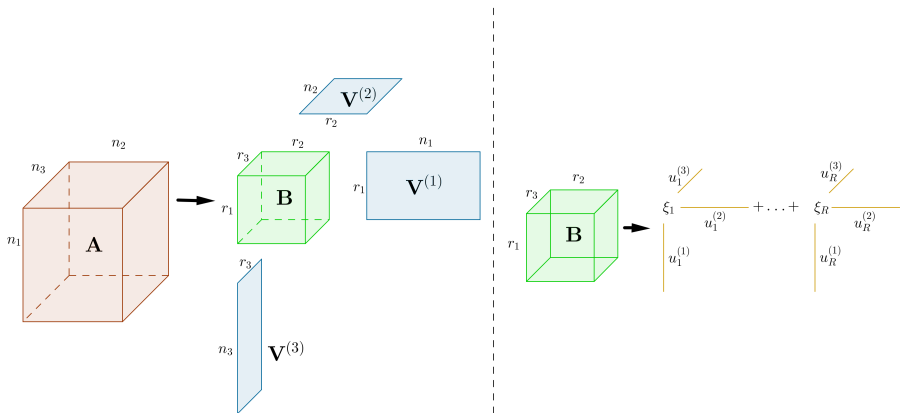


Figure: 3-d tensor  $\mathbf{A}$  in two-level Tucker-canonical format

# Multilinear algebra operations

Canonical format

$\mathbf{A}, \mathbf{B} \in \mathcal{C}_{R,n}$

$$\mathbf{A} = \sum_{k=1}^{R_A} c_k \mathbf{u}_k^{(1)} \otimes \dots \otimes \mathbf{u}_k^{(d)}, \quad \mathbf{B} = \sum_{j=1}^{R_B} d_j \mathbf{v}_j^{(1)} \otimes \dots \otimes \mathbf{v}_j^{(d)}. \quad (41)$$

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- scalar product:

$$\langle \mathbf{A}, \mathbf{B} \rangle = \sum_{k=1}^{R_A} \sum_{j=1}^{R_B} c_k d_j \prod_{\ell=1}^d \langle \mathbf{u}_k^{(\ell)}, \mathbf{v}_j^{(\ell)} \rangle \quad (42)$$

( $dR_A R_B n$  operations)



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- Hadamard product:

$$\mathbf{A} \odot \mathbf{B} = \sum_{k=1}^{R_A} \sum_{j=1}^{R_B} c_k d_j \bigotimes_{\ell=1}^d \left( \mathbf{u}_k^{(\ell)} \odot \mathbf{v}_j^{(\ell)} \right) \quad (43)$$

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# Multilinear algebra operations

## Canonical format

- discrete convolution:

$$\mathbf{A} * \mathbf{B} = \sum_{k=1}^{R_A} \sum_{j=1}^{R_B} c_k d_j \bigotimes_{\ell=1}^d (\mathbf{u}_k^{(\ell)} * \mathbf{v}_j^{(\ell)}) \quad (44)$$

( $dR_A R_B n \log n$  operations using 1D FFT)

# Multilinear algebra operations

Tucker format

$\mathbf{A}, \mathbf{B} \in \mathcal{T}_{r,n}$ :

$$\mathbf{A} = \mathbf{C} \times_1 \mathbf{U}^{(1)} \times_2 \dots \times_d \mathbf{U}^{(d)}, \quad \mathbf{B} = \mathbf{D} \times_1 \mathbf{V}^{(1)} \times_2 \dots \times_d \mathbf{V}^{(d)}. \quad (45)$$

# Multilinear algebra operations

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$(\mathcal{O}(dr^2n + r^{2d}))$  operations)

# Multilinear algebra operations

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# Multilinear algebra operations

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$(\mathcal{O}(dr^2 n \log n + r^{2d}))$  operations using 1D FFT)

# Multilinear algebra operations

## Summary

Tensor format	Scalar product	Hadamard product	Discrete convolution
Full	$n^d$	$n^d$	$n^d \log n$
Canonical	$dR_A R_B n$	$dR_A R_B n$	$dR_A R_B n \log n$
Tucker	$\mathcal{O}(dr^2 n + r^{2d})$	$\mathcal{O}(dr^2 n + r^{2d})$	$\mathcal{O}(dr^2 n \log n + r^{2d})$

Table: Computational complexities of various tensor formats

## 4. Low rank approximation of function related tensors



# Approximation of function related tensors

How to find an rank-structured approximation of tensor  $\mathbf{G} \in \mathbb{R}^{n^{\otimes d}}$  discretizing a multivariate function  $g : \mathbb{R}^d \rightarrow \mathbb{R}$  on a fine grid?

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- Tucker format
- Canonical format

# Approximation of function related tensors

## Tucker format

Consider a set  $\{\xi_j\}_{j=1}^{\mathbf{N}}$  of interpolation points, where  $\mathbf{N} = (N_1, \dots, N_d)$  represents a number of interpolation points at each dimension. The interpolant of the function  $g$ :

$$\mathcal{I}_{\mathbf{N}}g(\mathbf{x}) = \sum_{j=1}^{\mathbf{N}} g(\xi_j) l_j(\mathbf{x}), \quad l_j: \mathbb{R}^d \rightarrow \mathbb{R} \quad (49)$$

such as

$$l_j(\xi_j) = 1, \quad l_j(\xi_i) = 0 \quad i \neq j. \quad (50)$$

# Approximation of function related tensors

## Tucker format

Given orthonormal sets of interpolation functions

$$l_j = l_{j_1}^{(1)} \otimes \dots \otimes l_{j_d}^{(d)}, \quad (51)$$

we can write a Tucker approximation

$$g(\mathbf{x}) \approx \mathbf{G} = \Xi \times_1 \mathbf{I}^{(1)} \times_2 \dots \times_d \mathbf{I}^{(d)} \in \mathcal{T}_{\mathbf{N}, \mathbf{n}}, \quad (52)$$

where

$$[\Xi]_j = g(\xi_j) \quad (53)$$

and columns of side matrices  $\mathbf{I}^{(\ell)} \in \mathbb{R}^{n_\ell \times N_\ell}$  are discretized functions  $l_{j_\ell}^{(\ell)}$ ,  $j_\ell = 1, \dots, N_\ell$ .

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How to choose an orthonormal set of interpolating functions?

# Approximation of function related tensors

## Sinc function

The sinc function is defined as

$$\operatorname{sinc}(x) = \begin{cases} \frac{\sin(\pi x)}{\pi x}, & x \neq 0 \\ 1, & x = 0 \end{cases} \quad (54)$$

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Functions representing band-limited signals can be exactly reconstructed by their sampling values. So if the support of Fourier transform of given function,  $\hat{g}$ , is subset of  $[-\frac{\pi}{h}, \frac{\pi}{h}]$  ( $g \in U_h$ ), then

$$g(x) = \sum_{n=-\infty}^{\infty} g(nh) S_{n,h}(x), \quad (55)$$

where

$$S_{n,h}(x) = \operatorname{sinc}\left(\frac{x - nh}{h}\right). \quad (56)$$

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A set of functions  $\{S_{n,h}(x)\}_{n \in \mathbb{Z}}$  forms an orthonormal basis of  $U_h$ .



# Approximation of function related tensors

## Canonical format

Functions that depend on the sum of single variables, respectively their squares, i.e.

$$g(\mathbf{x}) = g(\rho), \quad \rho = \sum_{i=1}^d x_i, \quad \text{resp.} \quad \rho = \sum_{i=1}^d x_i^2. \quad (57)$$

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Assume that a function  $g(\rho)$  is given by transformation

$$g(\rho) = \int_{\Omega} G(t) e^{\rho F(t)} dt, \quad \Omega \in \{\mathbb{R}, \mathbb{R}_+, [a, b]\}. \quad (58)$$

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Suitable quadrature

$$g(\rho) \approx \sum_{j=1}^R \omega_j G(t_j) e^{\rho F(t_j)} \quad (59)$$

where  $\omega_j$  are quadrature heights,  $t_j$  are quadrature points.

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Fully separable representation!

$$g(\mathbf{x}) \approx \sum_{j=1}^R c_j \mu_j^{(1)} \otimes \dots \otimes \mu_j^{(d)}. \quad (61)$$

Sinc quadrature (rectangular rule):

$$\int_{\mathbb{R}} G(t) e^{\rho F(t)} dt \approx \int_{\mathbb{R}} \sum_{n=-M}^M S_{n,h}(t) G(nh) e^{\rho F(nh)} dt = \sum_{n=-M}^M G(nh) e^{\rho F(nh)} \quad (62)$$

# Approximation of function related tensors

Canonical approximation of newton kernel

$$f(\mathbf{x}) = \frac{1}{\|\mathbf{x}\|} = \frac{1}{\sqrt{x_1^2 + x_2^2 + x_3^2}}. \quad (63)$$

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After substitution:

$$\frac{1}{\rho} = \frac{2}{\sqrt{\pi}} \int_{\mathbb{R}} \cosh w \frac{e^{-\rho^2 \log^2(1+e^{\sinh w})}}{1 + e^{-\sinh w}} dw. \quad (66)$$

# Approximation of function related tensors

## Canonical approximation of newton kernel

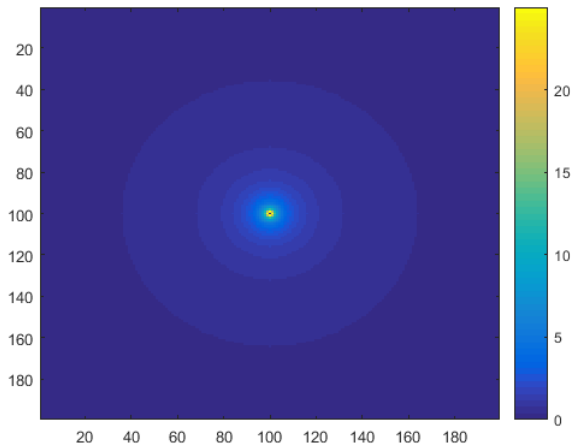


Figure: 2D newton potential discretized on square  $[-2, 2] \times [-2, 2]$

# Approximation of function related tensors

## Canonical approximation of newton kernel

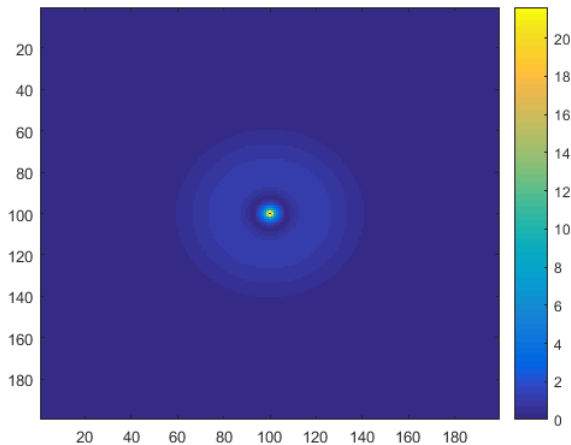


Figure: rank-3 approximation error

# Approximation of function related tensors

## Canonical approximation of newton kernel

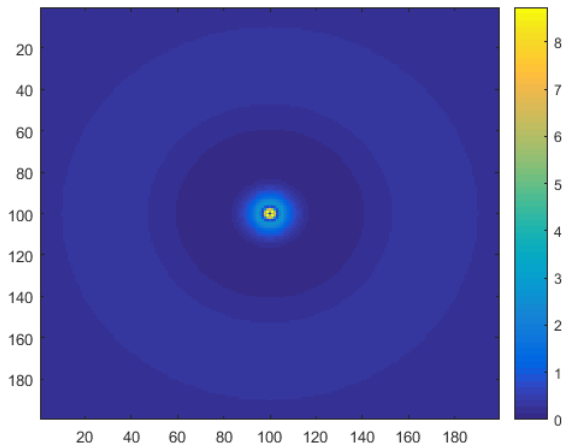


Figure: rank-5 approximation error

# Approximation of function related tensors

## Canonical approximation of newton kernel

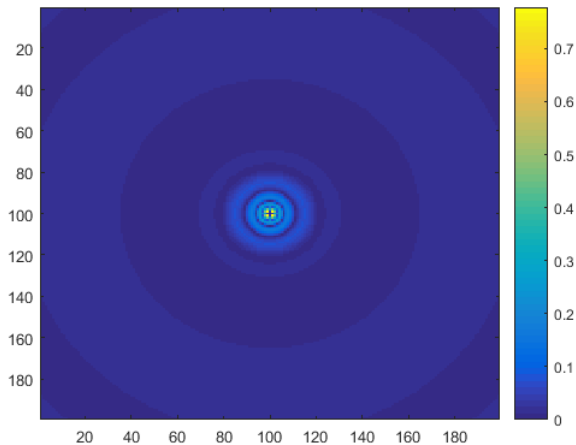


Figure: rank-11 approximation error

# Approximation of function related tensors

## Canonical approximation of newton kernel

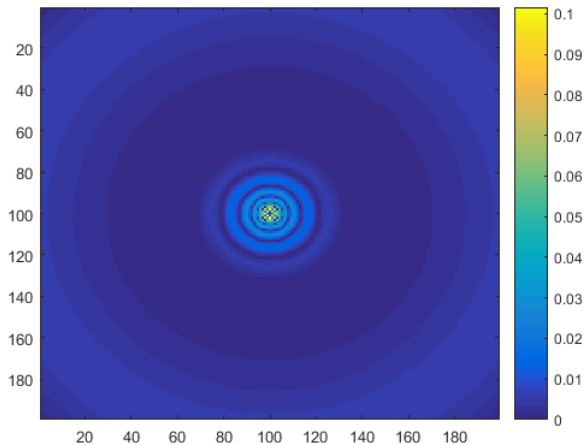


Figure: rank-15 approximation error

# Other applications



- tensor representation of multidimensional operators

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- superfast FFT

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**Thank you for attention!**

- Boris Khoromskij homepage:  
<http://personal-homepages.mis.mpg.de/bokh/>
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